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Subject: Re: Randomize array order

Posted by [Conor](#) on Fri, 27 Jul 2007 16:44:53 GMT

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On Jul 27, 12:05 pm, kuyper <kuy...@wizard.net> wrote:

> David Streutker wrote:

>> On Jul 27, 6:03 am, Allan Whiteford

> ...

>>> If you have a million elements then you have 1000000! (i.e. one million  
>>> factorial) different ways to re-order the data. However, your seed is a  
>>> 4 byte integer which can only take  $2^{32}$  different values.

>

>>> Some messing about suggests that:

>

>>>  $1000000! \approx 10^{5568636}$

>

>>> which means there are  $\sim 10^{5568636}$  different ways to re-arrange your  
>>> elements as opposed to the  $4 \times 10^9$  values your seed can take.

>

>>> Thus, using any of the algorithms suggested you're only going to sample

>

>>>  $10^{-5568625} \%$

>

>>> of the possible values. This is a really small number. It means that no  
>>> matter how hard you try and how many times you do things you'll never be  
>>> able to access anything but a tiny number of the possibilities without  
>>> doing multiple shufflings - I think it's something like 618737  
>>> sub-shufflings (i.e.  $5568636 / 9$ ) but that could be wrong. However, that  
>>> requires producing 618737 seeds per major-shuffle (and you can't use a  
>>> generator based on a 4 byte seed to produce these seeds).

>

>>> But, since you're only going to be running the code 1000-10,000 times  
>>> (which is much smaller than  $4e9$ ) I guess everything will be ok. I don't  
>>> know if anyone has studied possible correlations of results as a  
>>> function of the very small number of seeds (compared to the data),  
>>> whatever random number generator is used and the shuffling method.  
>>> Presumably they have and presumably everything is ok. Does anyone know?

>

>>> Thanks,

>

>>> Allan

>

>> I'm not sure that I agree. Where in any of our algorithms are we  
>> unable to access a (theoretically) possible outcome? As long as we  
>> are able to randomly select any element of the array in each step, it  
>> should work, right? (i.e., as long as the input array has fewer than  
>>  $2^{32}$  elements.) In your analysis, shouldn't we be using  $(2^{32})^n$  for  
>> the maximum possible number of randomly generated combinations, where

>> n is the number of steps/elements?

>

> No, because the entire sequence of numbers is uniquely determined by  
> initial internal state of the generator. If you knew the algorithm  
> used, and the internal state, that's all the information you'd need to  
> predict, precisely, the entire sequence of numbers generated, no  
> matter how long that sequence was. If the internal state is stored in  
> a 32 bit integer, that means there's only  $2^{32}$  possible different  
> sequences.

>

>> From that fact, it can also be shown that every possible sequence must

>

> start repeating, exactly, with a period that is less than  $2^{32}$ . If one  
> of the possible sequences has starts repeating with a period T, then  
> at least T-1 of the other possible sequences generate that same repeat  
> cycle, with various shifts.

>

> There's a reason why these things are called PSEUDO-random number  
> generators.

It shouldn't really be a problem for me, fortunately. I'm running  
this a couple thousand times, but everytime it is on a different set  
of values. The only thing I would have to worry about is it repeating  
within one set of values, which won't happen for 1,000,000 elements.

Of course, worse comes to worse there's always a true random number  
generator:

[www.random.org](http://www.random.org)

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