
Subject: Re: IDL interpreter questions - can someone (D.Fanning) explain - TIA
Posted by [Craig Markwardt](#) on Fri, 18 May 2001 19:31:36 GMT
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<mankoff@I.HATE.SPAM.lasp.colorado.edu> writes:

> On Fri, 18 May 2001, JD Smith wrote:
>
>> dadada wrote:
>>> How are variables referenced by default?
>> I'm not sure what you mean here. Pointer references? They are explicit
>> only... i.e. you can't create a reference of an existing variable.
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> Not sure either, but here is my interpretation of the question/answer:
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> In functions, variables are *always* 'by value'
> In procedures, they are 'by value' unless you put a "return" statement
> anywhere in the procedure. If this exists, then they are passed 'by
> reference'

Sorry Ken I'm going to have to take you to task for a few things.
First of all, pass by value vs. pass by reference:

- * all variables are passed by reference, *except*
- * subscripted arrays, structure tags, and (I believe) system variables,
which are passed by value

It doesn't make a difference whether you have a return statement or not.

As for continuations, closures, etc., these are computer science jargon for specific language behaviors. IDL has none of them. I understand continuations to be a way for execution contexts to be suspended, saved, and later restored. Perhaps the CATCH error handling technique is a nascent continuation. Alas, this has nothing to do with the CONTINUE reserved word recently added to FOR and WHILE loops.

Craig

--

Craig B. Markwardt, Ph.D. EMAIL: craigmnet@cow.physics.wisc.edu
Astrophysics, IDL, Finance, Derivatives | Remove "net" for better response

Subject: Re: IDL interpreter questions - can someone (D.Fanning) explain - TIA
Posted by [<mankoff@I.HATE.SPAM.lasp.colorado.edu>](#) on Sat, 19 May 2001 00:17:33 GMT
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On 18 May 2001, Craig Markwardt wrote:

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> It doesn't make a difference whether you have a return statement or
> not.
```

Hi JD,

No need to apologize for a correction. But i have some questions about this that maybe you can answer.

I thought 'by value' meant that the called routine gets a copy of the variable, and cannot modify the contents of the variable in the calling routine.

And that 'by reference' means that the called routine gets a pointer to the variable from the calling routine. Any changes that the called makes, appear in the caller.

Now its true that I don't know anything about the actual IDL implementation (though I have written RPC code for IDL). I actually answered based upon the behavior of IDL, not the implementation. That is, functions won't modify the callers variables, and neither will procedures, unless you add the 'return'.

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- > suspended, saved, and later restored. Perhaps the CATCH error
- > handling technique is a nascent continuation. Alas, this has nothing
- > to do with the CONTINUE reserved word recently added to FOR and WHILE
- > loops.

ok, i get it.

But if you want to "suspend, save, restore" the execution state, wouldn't "save, /vars" and "save, /all" simulate this to some extent?

thanks for any clarifications,

-k.

--

Ken Mankoff

LASP://303.492.3264

<http://lasp.colorado.edu/~mankoff/>

Subject: Re: IDL interpreter questions - can someone (D.Fanning) explain - TIA
Posted by [Mark Rivers](#) on Sat, 19 May 2001 05:28:00 GMT
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mankoff@I.HATE.SPAM.lasp.colorado.edu wrote in message ...

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- > answered based upon the behavior of IDL, not the implementation. That is,
- > functions won't modify the callers variables, and neither will procedures,
- > unless you add the 'return'.

That's not true. Here's the proof:

```
pro myproc, a
```

```
    a=2
```

```
end
```

```
function myfunc, a
```

```
    a=3
```

```
end
```

```
; Test program
```

```
a = 1
```

```
myproc, a
```

```
help, a
```

```
a = 0
```

```
t = myfunc(a)
help, a
end
```

```
.run test
```

```
A      INT    =    2
A      INT    =    3
```

So the procedure and the function both modified arguments passed to them.

Mark Rivers

Subject: Re: IDL interpreter questions - can someone (D.Fanning) explain - TIA
Posted by [Craig Markwardt](#) on Sat, 19 May 2001 17:20:12 GMT
[View Forum Message](#) <> [Reply to Message](#)

<mankoff@I.HATE.SPAM.lasp.colorado.edu> writes:

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> ok, i get it.
>
> But if you want to "suspend, save, restore" the execution state, wouldn't
> "save, /vars" and "save, /all" simulate this to some extent?
```

Yes, sort of, but a formal continuation also saves the execution context, not just the variables. Take a look at CATCH for exception handling and you will see that CATCH can essentially return *twice*. Once in the normal program flow, and a second time if an error occurs.

So, internally CATCH saves the execution context so that when an error happens the execution can resume again within CATCH. To be honest I don't totally understand formal continuations, but I think CATCH is a rudimentary example of one.

Craig

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Subject: Re: IDL interpreter questions - can someone (D.Fanning) explain - TIA
Posted by [John-David T. Smith](#) on Tue, 22 May 2001 20:41:25 GMT
[View Forum Message](#) <> [Reply to Message](#)

mankoff@I.HATE.SPAM.lasp.colorado.edu wrote:

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> this that maybe you can answer.
```

I plead the 5th. Craig, I believe, was the taskmaster. I *never* take
anyone to task <sarcasmicon>.

JD
